

SPSC: a Simple Supercompiler in Scala

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PU'09

To program is to understand



Kristen Nygaard

A program should reflect
an understanding of
the problem domain



We would like...

- To demystify supercompilation for a programmer
- In order to do this we want:
 - To present the core of supercompilation in the form of a program
 - The code of the supercompiler should be small and clear
 - A minimalistic input language

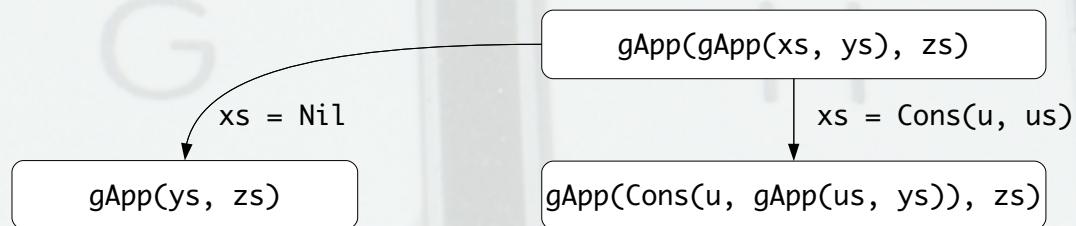
Supercompilation in a Nutshell

```
gApp(Nil(), vs) = vs;  
gApp(Cons(u, us), vs) = Cons(u, gApp(us, vs));
```

```
gApp(gApp(xs, ys), zs)
```

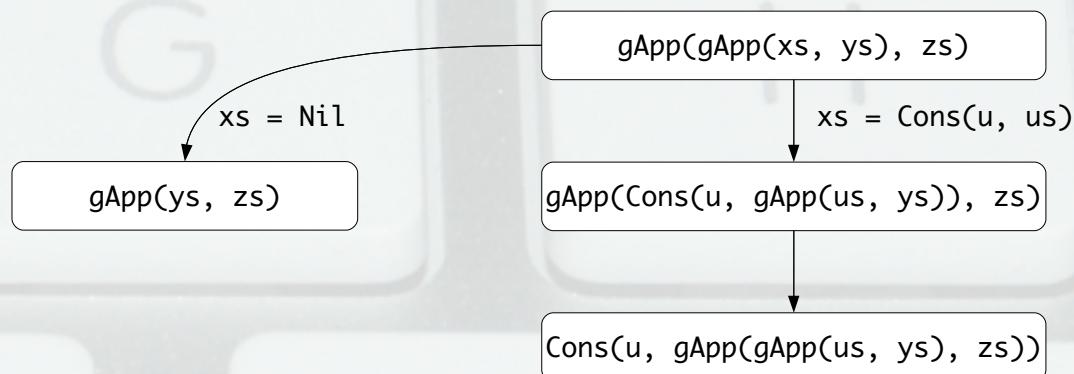
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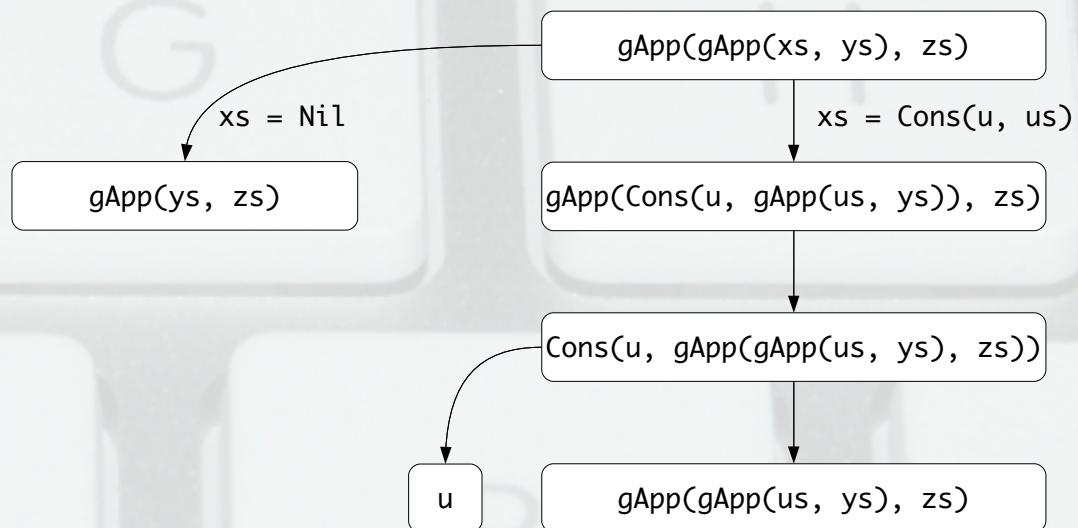
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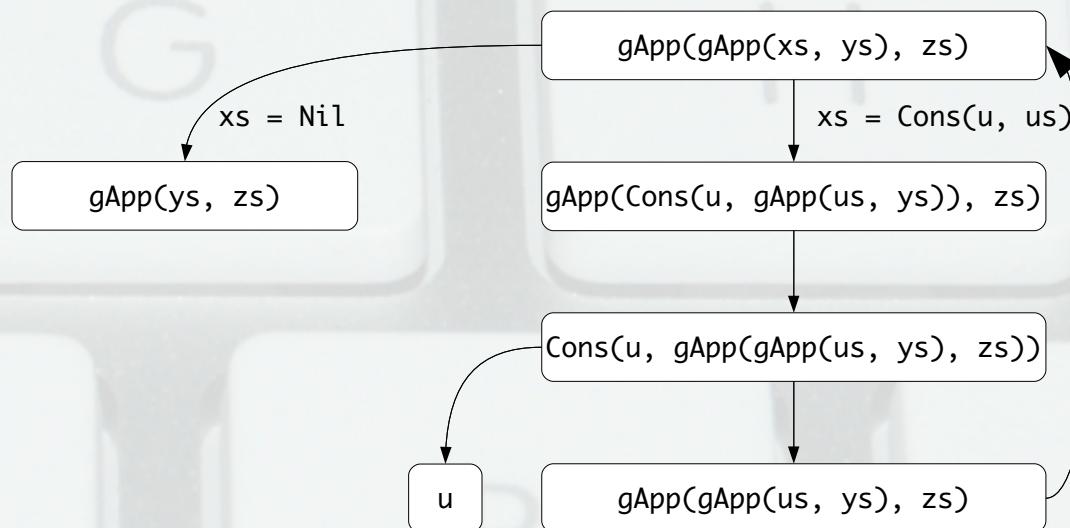
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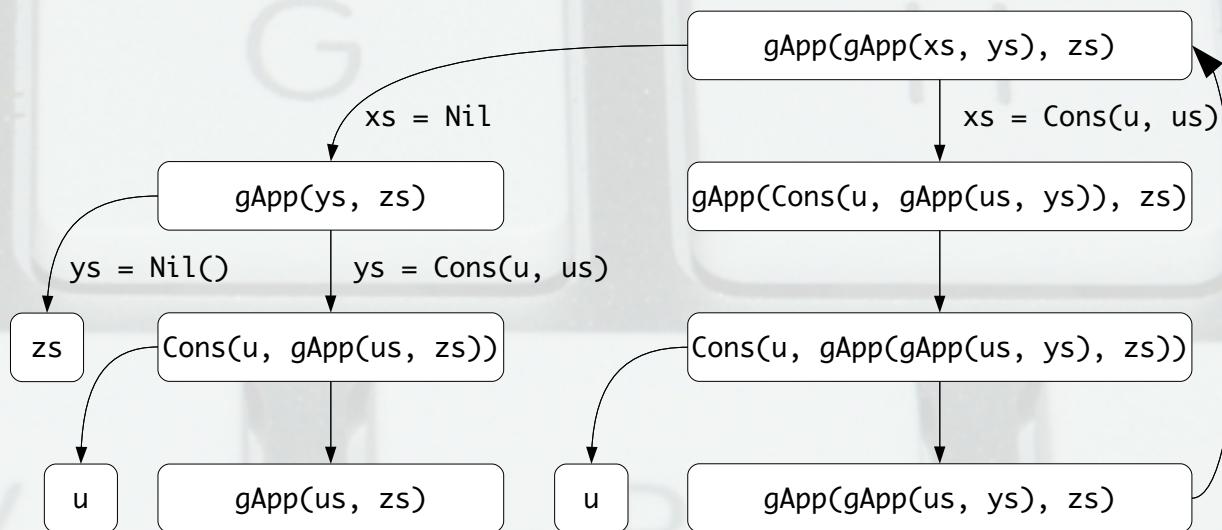
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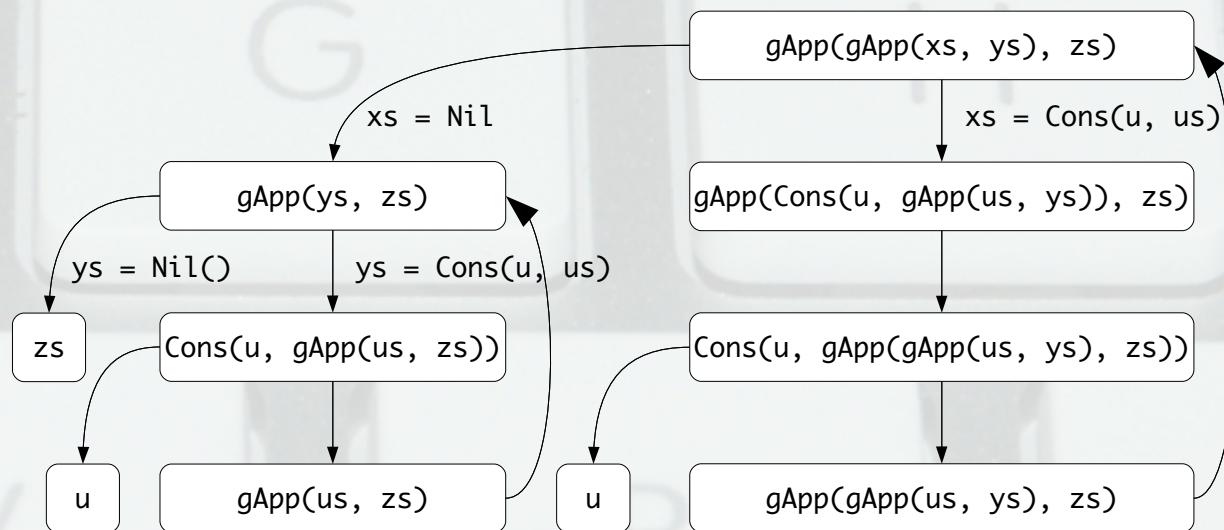
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Supercompilation in a Nutshell

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gApp(Nil(), vs)      = vs;  
gApp(Cons(u, us), vs) = Cons(u, gApp(us, vs));
```



Supercompilation in a Nutshell

```
gApp(Nil(), vs)      = vs;  
gApp(Cons(u, us), vs) = Cons(u, gApp(us, vs));
```

gApp(gApp(xs, ys), zs)

gApp1(xs, ys, zs)

```
gApp1(Nil(), y, z)      = gApp2(y, z);  
gApp1(Cons(v1, v2), y, z) = Cons(v1, gApp1(v2, y, z));
```

```
gApp2(Nil(), z)          = z;  
gApp2(Cons(v3, v4), z)    = Cons(v3, gApp2(v4, z));
```

Quest for practical supercompilers

- 1974 V. Turchin presented supercompilation to a group of students at seminars in Moscow
- 1980s V. Turchin developed first supercompilers for the functional programming language Refal (CUNY, New York)
- 1980s – Papers by V. Turchin on supercompilation of Refal
 - 1990s
- 1993 – Andrei Nemytykh (IPS RAS, Pereslavl-Zalesky) continued work on
 - 2000s Turchin's supercompiler and completed it
- 1998 – Java Supercompiler by Andrei Klimov, Arkady Klimov and Artem Shvorin
 - 2000s (Keldysh Institute of Applied Mathematics, RAS, Moscow)

[AK]

Quest for understandable supercompilers

- 1986 V. Turchin "The Concept of a Supercompiler".
ACM TOPLAS 8(3): 292-325
- 1988 V. Turchin "The Algorithm of Generalization in the Supercompiler".
Partial Evaluation and Mixed Computation, North-Holland, 341-353
- 1993 ...
The first understandable supercompiler and paper by R. Glück and
A. Klimov "Occam's razor in metacomputation: the notion of a perfect
process tree". LNCS 724: 112-123
- 1993 Book by S. Abramov "Metacomputation and its application" (in Rus)
- 1996 M.H. Sørensen, R. Glück, N.D. Jones "A Positive Supercompiler". J.
Funct. Program. 6(6): 811-838
- 2009 I. Klyuchnikov, S. Romanenko "SPSC: a Simple Supercompiler in Scala".
PU'09

[AK]

Input Language

$p ::= d_1 \dots d_n \text{ program}$

$d ::= f(x_1, \dots, x_n) = e; \quad \text{f-function}$

$\quad | \quad g(q_1, x_1, \dots, x_n) = e_1; \quad \text{g-function}$

\dots

$\quad g(q_m, x_1, \dots, x_n) = e_m;$

$e ::= x \text{ variable}$

$\quad | \quad c(e_1, \dots, e_n) \quad \text{constructor}$

$\quad | \quad f(e_1, \dots, e_n) \quad \text{call to f-function}$

$\quad | \quad g(e_1, \dots, e_n) \quad \text{call to g-function}$

$q ::= c v_1 \dots v_n \quad \text{pattern}$

Implementation language??

- Easy to understand
- Easy to use (IDE, debugger, libs, ...)
- Functional
- Also cool

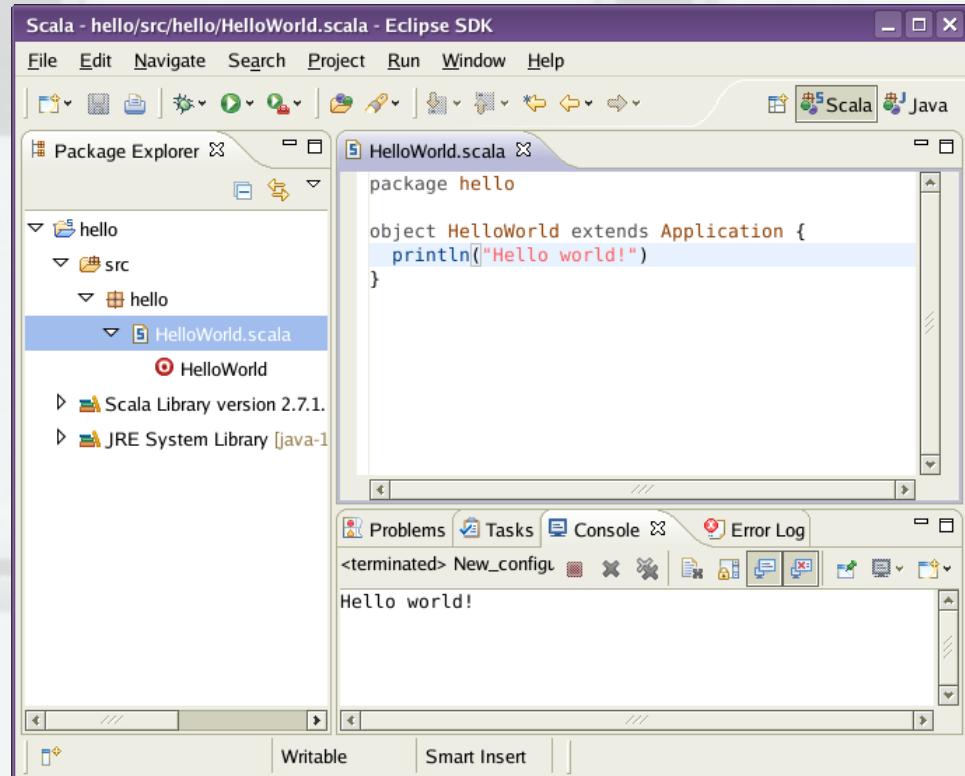
Scala makes a buzz



Scalable Language

- FP support
- OOP support
- Compiles to JVM bytecode
- Ready for production – Twitter is rewritten into Scala

Scala makes a buzz



- IDE Support
- Convenient debugger
- Great community
- Strong theoretical base

Scala by Example

pragmatic OOP

```
abstract class Def {def name: String}

  class FFun(name: String, args: List[Var], term: Term) extends Def {
    override def toString =
      name + args.mkString("(, ", " ,")"") + " = " + term + ";"
  }

  class GFun(name: String, p: Pat, args: List[Var], term: Term) extends Def {
    override def toString =
      name + (p :: args).mkString("(, ", " ,")"") + " = " + term + ";"
  }
```

Scala by Example

pattern matching with case classes

```
abstract class Def {def name: String}

case class FFun(name: String, args: List[Var], term: Term) extends Def {
  override def toString =
    name + args.mkString("(, ", " ,")"") + " = " + term + ";"
}

case class GFun(name: String, p: Pat, args: List[Var], term: Term) extends Def {
  override def toString =
    name + (p :: args).mkString("(, ", " ,")"") + " = " + term + ";"
}

case class Program(defs: List[Def]){
  val f = (defs :\ Map[String, FFun]())
  {case (x: FFun, m) => m + (x.name -> x); case (_, m) => m}
  val g = (defs :\ Map[(String, String), GFun]())
  {case (x: GFun, m) => m + ((x.name, x.p.name) -> x); case (_, m) => m}
  val gs = (defs :\ Map[String, List[GFun]]() .withDefaultValue(Nil))
  {case (x: GFun, m) => m + (x.name -> (x :: m(x.name))); case (_, m) => m}
  override def toString = defs.mkString("\n")
}
```

Scala by Example

functional objects

```
abstract class Def {def name: String}

case class FFun(name: String, args: List[Var], term: Term) extends Def {
  override def toString =
    name + args.mkString("(, ", " ,")"") + " = " + term + ";"
}

case class GFun(name: String, p: Pat, args: List[Var], term: Term) extends Def {
  override def toString =
    name + (p :: args).mkString("(, ", " ,")"") + " = " + term + ";"
}

case class Program(defs: List[Def]){
  val f = (defs :\ Map[String, FFun]())
  {case (x: FFun, m) => m + (x.name -> x); case (_, m) => m}
  val g = (defs :\ Map[(String, String), GFun]())
  {case (x: GFun, m) => m + ((x.name, x.p.name) -> x); case (_, m) => m}
  val gs = (defs :\ Map[String, List[GFun]]() .withDefaultValue(Nil))
  {case (x: GFun, m) => m + (x.name -> (x :: m(x.name))); case (_, m) => m}
  override def toString = defs.mkString("\n")
}
```

```
val g = p.g(name, cname)
```

Scala by Example

higher-order functions, almost any names for defs...

```
abstract class Def {def name: String}

case class FFun(name: String, args: List[Var], term: Term) extends Def {
  override def toString =
    name + args.mkString("(, ", " ,")"") + " = " + term + ";"
}

case class GFun(name: String, p: Pat, args: List[Var], term: Term) extends Def {
  override def toString =
    name + (p :: args).mkString("(, ", " ,")"") + " = " + term + ";"
}

case class Program(defs: List[Def]){
  val f = (defs :\ Map[String, FFun]())
    {case (x: FFun, m) => m + (x.name -> x); case (_, m) => m}
  val g = (defs :\ Map[(String, String), GFun]())
    {case (x: GFun, m) => m + ((x.name, x.p.name) -> x); case (_, m) => m}
  val gs = (defs :\ Map[String, List[GFun]]() .withDefaultValue(Nil))
    {case (x: GFun, m) => m + (x.name -> (x :: m(x.name))); case (_, m) => m}
  override def toString = defs.mkString("\n")
}
```

Scala by Example

to mutate (and how) or not to mutate?

```
def findSubst(t1: Term, t2: Term) = {
    val map = scala.collection.mutable.Map[Var, Term]()
    def walk(t1: Term, t2: Term): Boolean = (t1, t2) match {
        case (v1: Var, _) => map.put(v1, t2) match {
            case None => true;
            case Some(t3) => t2 == t3
        }
        case (e1: CFG, e2: CFG) if shellEq(e1, e2) =>
            List.forall2(e1.args, e2.args)(walk)
        case _ => false
    }
    if (walk(t1, t2)) map.readOnly else null
}
```

Scala by Example

to mutate (and how) or not to mutate?

```
def findSubst(t1: Term, t2: Term) = {  
    var map = Map[Var, Term]()  
    def walk(t1: Term, t2: Term): Boolean = (t1, t2) match {  
        case (v1: Var, _) => map.get(v1) match {  
            case None => map += (v1 -> t2); true  
            case Some(t3) => t2 == t3  
        }  
        case (e1: CFG, e2:CFG) if shellEq(e1, e2) =>  
            List.forall2(e1.args, e2.args)(walk)  
        case _ => false  
    }  
    if (walk(t1, t2)) map else null  
}
```

Scala by Example

real DSL

```
object SParsers extends StandardTokenParsers with ImplicitConversions {  
    ...  
    def prog = definition+  
    def definition: Parser[Def] = gFun | fFun  
    def term: Parser[Term] = fcall | gcall | ctr | vrb  
    def vrb = lid ^^ Var  
    def pat = uid ~ ("(" ~> repsep(vrb, ",") <~ ")") ^^ Pat  
    def fFun =  
        fid ~ ("(" ~> repsep(vrb, ",") <~ ")") ~ ("=" ~> term <~ ";") ^^ FFun  
    def gFun =  
        gid ~ ("(~>pat) ~ (((", "~>vrb)*)) ~ ("=" ~> term <~ ";") ^^ GFun  
    def ctr = uid ~ ("(" ~> repsep(term, ",") <~ ")") ^^ Ctr  
    def fcall = fid ~ ("(" ~> repsep(term, ",") <~ ")") ^^ FCall  
    def gcall = gid ~ ("(" ~> repsep(term, ",") <~ ")") ^^ GCall  
}
```

Time to try it

spsc_web beta

http://localhost:8080/spsc_web/spsc_result

spsc_web beta

Input code

```
append(Nil, vs) = vs;
append(Cons(u, us), vs) = Cons(u, append(us, vs));
appendXYaZ(xs, ys, zs) = append(append(xs, ys), zs);
```

Function to be supercompiled

```
appendXYaZ
```

Supercompiled code

```
append1(Nil(), a) = a;
append1(Cons(a, b), c) = Cons(a, append1(b, c));
append2(Nil(), a) = a;
append2(Cons(a, b), c) = Cons(a, append2(b, c));
appendXYaZ(a, b, c) = append1(append2(a, b), c);
```

Partial process tree

```
graph TD
    Root[appendXYaZ(xs, ys, zs)] --> Let[let $248=append(xs, ys), $249=zs in append($248, $249)]
    Let --> Append1[append($248, $249)]
    Let --> Zs[zs]
    Append1 --> Cons1[$248=Cons($256, $257)]
    Append1 --> Nil1[$248=Nil()]
    Append1 --> Cons2[xs=Cons($250, $251)]
    Append1 --> Nil2[xs=Nil()]
```

Thanks!

Writing it is easy,
understanding it is hard.

Anonymous

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